Asymptotic properties of graphs on orientable surfaces

Mihyun Kang





Doctoral School Discrete Mathematics

Outline

I. Erdős-Rényi Random Graphs

II. Random Graphs on Surfaces

III. Proof Ideas

Part I. Erdős–Rényi Random Graph *G*(*n*, *m*)

chosen uniformly at random among all graphs with n vertices and m edges



Emergence of "Giant" Component

L(d) = # vertices in the largest component in G(n, m), where $m = d \cdot \frac{n}{2}$

Theorem		[ERDŐS-RÉNYI 60]
● If <i>d</i> < 1, whp	$L(d) = O(\log n)$	
• If <i>d</i> = 1, whp	$L(d) = \Theta(n^{2/3})$	
• If <i>d</i> > 1, whp	$L(d) = \Theta(n)$	
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•••		
•		
• •	••	

 $\Theta(n^{2/3})$

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Θ(n)

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• If d < 1, whp $L(d) = O(\log n)$

and the largest component is a tree and all other components are either trees or unicyclic components.

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- If d = 1, whp $L(d) = \Theta(n^{2/3})$
- If d > 1, whp $L(d) = (1 + o(1)) \rho n$

where $\rho = \rho(d) \in (0, 1)$ is the unique positive solution of $1 - \rho = e^{-d\rho}$,

and the largest component contains more than two cycles, while all but the largest component are trees or unicyclic components of order $O(\log n)$.

Critical Phenomenon

Let $d = 1 + \epsilon$ for $\epsilon = o(1)$.

Theorem

[BOLLOBÁS 84; ŁUCZAK 90]

- If $\epsilon n^{1/3} \to -\infty$, whp
- If $\epsilon n^{1/3} \rightarrow \lambda$, a constant, whp
- If $\epsilon n^{1/3} \to +\infty$, whp

$$L(1 + \epsilon) = 2\epsilon^{-2} \log |\epsilon|^3 n \ll n^{2/3}$$
$$L(1 + \epsilon) = \Theta(n^{2/3})$$
$$L(1 + \epsilon) = (1 + o(1)) 2\epsilon n \gg n^{2/3}$$



Part II. Random Graphs on Surfaces



Mihyun Kang (TU Graz, Austria) Graphs on orientable surfaces

Embeddability on \mathbb{S}_g

Let \mathbb{S}_g be the orientable surface of genus g > 0.

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Facewidth

- The facewidth of a map is the minimum number of intersections that a non-contractible circle has with the map.
- The facewidth of a graph is the maximal facewidth of all its embeddings.

Random Planar Graphs

Let P(n, m) be a uniform random planar graph with *n* vertices and *m* edges and L(d) denote the number of vertices in the largest component in

 $P(n, d \cdot \frac{n}{2})$. Let $d = 1 + \epsilon$ for $\epsilon = o(1)$.

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• If $\epsilon n^{2/5} \to -\infty$, whp If $\epsilon n^{2/5} \to \infty$, whp	$n - L(2 + \epsilon) = \epsilon n \gg n^{3/5}$ $n - L(2 + \epsilon) = \Theta(\epsilon^{-3/2}) \ll n^{3/5}$

 $1 >> n^{2/3}$

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Random Graphs with Genus g > 0

Let $S_g(n, m)$ be a uniform random graph on a surface of genus g > 0with *n* vertices and *m* edges and L(d) denote the number of vertices in the largest component in $S_g(n, d \cdot \frac{n}{2})$. Let $d = 1 + \epsilon$ for $\epsilon = o(1)$.

Theorem	[K Mosshammer-Sprüssel 15+]
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Mihyun Kang (TU Graz, Austria) Graphs on orientable surfaces

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 $n - L << n^{3/5}$

Part III. Proof Ideas

- Internal structure: Core-Kernel
- Constructive decomposition along connectivity
- Singularity analysis of generating functions
- Probabilistic analysis



complex com. unicyc. com. trees



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- \Rightarrow Core of G: maximal subgraph of G with minimum degree two
- \Rightarrow Kernel of G: obtained from core of G by replacing each path by an edge
- \Rightarrow G is embeddable on \mathbb{S}_g if and only if kernel of G is embeddable on \mathbb{S}_g



▷ Cubic weighted multigraphs

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- \triangleright Constructive decomposition: 1-conn. \implies 2-conn.

Theorem

[ROBERTSON-VITRAY 90]

• Every 1-connected graph *C* embeddable on S_g with facewidth $k \ge 2$ has a unique 2-conn. component embeddable on S_g with facewidth *k*, while all other 2-connected components are planar.

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- Every 3-connected graph embeddable on S_g has a unique embedding on S_g if its facewidth is at least 2g + 3

- Cubic weighted multigraphs
- \triangleright Constructive decomposition: 1-conn. \Longrightarrow 2-conn. \Longrightarrow 3-conn.
- \triangleright Dual of 3-connected cubic maps on \mathbb{S}_g are triangulations on \mathbb{S}_g

in which separating double edges, separating loops, and separating pair of loops are forbidden

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Generating Functions

Theorem

$$G_g(y) = \sum_{\sum g_i \le g} \frac{1}{k!} \prod_{i=1}^k C_{g_i}(y)$$

$$C_g(y) \sim C_g^{\text{fw} \ge 2}(y) \sim B_g^{\text{fw} \ge 2} \left(\frac{y}{1 - Q_0(y)}\right)$$

$$B_g(y) \sim B_g^{fw \ge 3}(y) \sim T_g^{fw \ge 3}(y(1 + N_0(y)))$$

$$T_g^{\text{fw}\geq 3}(y) \sim T_g^{\text{fw}\geq 2g+3}(y) \sim S_g(y)$$

- $G_g(y)$: cubic multigraphs on \mathbb{S}_g
- $C_q(y)$: 1-connected cubic multigraphs on \mathbb{S}_q
- $Q_0(y)$: 1-connected cubic multigraphs on S_0 rooted at a loop
- $B_g(y)$: 2-connected cubic multigraphs on \mathbb{S}_g
- $N_0(y)$: 2-connected cubic multigraphs on \mathbb{S}_0 rooted at an edge
- $T_q(y)$: 3-connected cubic graphs on \mathbb{S}_q
- $S_g(y)$: simple triangulations on \mathbb{S}_g

Theorem

[FANG-K.- MOSSHAMMER-SPRÜSSEL 15]

Let $K_g(n)$ denote the number of cubic weighted multigraphs embeddable on \mathbb{S}_g with *n* vertices, for $n \in \mathbb{N}$ even and $g \ge 0$. Then

$$K_g(n) \sim c_g n^{5/2(g-1)-1} \gamma^n n!$$

where c_g is a constant depending only on genus g and $\gamma = \frac{79^{3/4}}{54^{1/2}} \approx 3.6$.

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• \mathcal{G} is not embeddable on \mathbb{S}_{g-1} ;

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- \mathcal{G} is not embeddable on \mathbb{S}_{g-1} ;
- the largest component in *G* is a unique non-planar component and is of order *n* − *O*(1).

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Construct complex graph *G* embeddable on S_g by

• choosing the kernel of G from the set of possible candidates

$$=\sum_{i,d} K_g(2\ell-d)$$

Construct complex graph *G* embeddable on S_g by

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- putting on its edges vertices of degree two to obtain the core of G

$$=\sum_{i,d} \mathcal{K}_g(2\ell-d) \frac{(k)_i}{(2\ell-d)!} \begin{pmatrix} i-a\ell-1\\ 3\ell-d-1 \end{pmatrix}$$

Construct complex graph G embeddable on \mathbb{S}_g by

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- adding a forest rooted at vertices of the core of G

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 $C_g(k, k + \ell) = \#$ complex graphs on \mathbb{S}_g with k vertices and $k + \ell$ edges

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in which main contribution comes from the terms

$$i = (1 + O(\sqrt{\ell/k}) + O(1/\sqrt{\ell}))\sqrt{3k\ell}$$
 and $d = \Theta(\sqrt{\ell^3/k})$

 $S_g(n,m) = \#$ graphs on \mathbb{S}_g with *n* vertices and *m* edges

 $C_g(k, k + \ell) = \#$ complex graphs on \mathbb{S}_g with k vertices and $k + \ell$ edges $U(n - k, m - k - \ell) = \#$ graphs without complex components with n - k vertices and $m - k - \ell$ edges

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$$k = (1 + o(1)) \varepsilon n$$
 and $\ell = (\alpha + o(1)) \varepsilon n^{1/3}$.

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• When $m = (2 + \varepsilon)\frac{n}{2}$, $n^{-2/5} \ll \varepsilon \ll 1$, the main contribution comes from

$$k = n - (\beta + o(1)) \varepsilon^{-3/2}$$
 and $\ell = t + (2\beta + o(1)) \varepsilon^{-3/2}$

Two Critical Periods

Let $S_g(n, m)$ be a graph chosen uniformly at random among all graphs embeddable on \mathbb{S}_g of genus $g \ge 1$ with *n* vertices and *m* edges.

Let C_1 denote the largest component in $S_g(n, m)$.

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$$|n-|\mathcal{C}_1| = \Theta(\varepsilon^{-3/2}) \ll n^{3/5}$$

Overview

G(n, m) Erdős–Rényi random graph

[ERDŐS-RÉNYI 60]

- Phase transition and critical phenomenon
- Emergence of giant component when # edges $m = \frac{n}{2} + O(n^{3/2})$



Overview



Outlook

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Conjecture

- There exist $\alpha = \alpha(n), \beta = \beta(n)$ satisfying $\alpha(n) \le \beta(n)$ such that
 - if g ≪ α,
 S_g(n, m) exhibits the second critical phase analogous to P(n, m), but with different critical exponent
 - if $g \gg \beta$,

 $S_g(n, m)$ does not exhibit the second critical phase analogous to G(n, m)