

Phase Transitions in Random Hypergraphs

Mihyun Kang

Joint work with Oliver Cooley and Christoph Koch



Emergence of Giant Component in $G(n, p)$

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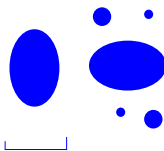
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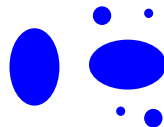
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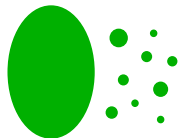
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Let $d = p \cdot (n - 1) = 1 + \epsilon$ for $\epsilon = o(1)$.

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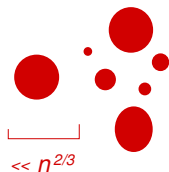
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[BOLLOBÁS 84; ŁUCZAK 90; ALDOUS 97]

• If $\epsilon^3 n \rightarrow -\infty$, whp

$$L_1(d) \sim 2\epsilon^{-2} \log \epsilon^3 n \ll n^{2/3}$$



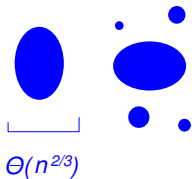
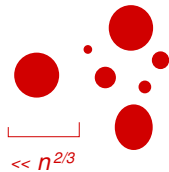
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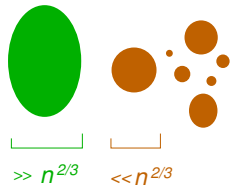
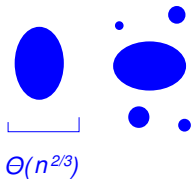
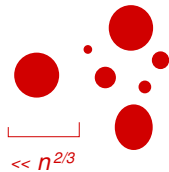
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- If $\epsilon^3 n \rightarrow +\infty$, whp $L_1(d) \sim 2\epsilon n \gg n^{2/3}$



Asymptotic Normality of Giant Component

Assume $d = p \cdot (n - 1) > 1$ and $0 < \rho < 1$ satisfies $1 - \rho = e^{-d \cdot \rho}$.

Let $\mu := \rho \cdot n$ and $\sigma^2 := \frac{\rho(1-\rho)}{(1-d(1-\rho))^2} \cdot n$

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Central limit theorem

Let $N(0, 1)$ denote the standard normal distribution. Then

$$\frac{L_1(d) - \mu}{\sigma} \xrightarrow{d} N(0, 1)$$

for d constant

[STEPANOV 70; BEHRISCH-COJA-OGHLAN-K. 09]

for $(d - 1)^3 n \rightarrow \infty$

[PITTEL-WORMALD 05; BOLLOBÁS-RIORDAN 12]

Proof techniques

- Counting connected graphs inside-out [PW 05]
- Stein's method [BC-OK 09]
- Random walk [BR 12]

Local Limit Theorem for Giant Component

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[STEPANOV 70; PITTEL-WORMALD 05; BEHRISCH-COJA-OGHLAN-K. 09]

Let $d > 1$ be constant and $I \subset \mathbb{R}$ compact. For any $k \in \mathbb{N}$ with $\sigma^{-1}(k - \mu) \in I$

$$\mathbb{P}[L_1(d) = k] \sim \frac{1}{\sigma\sqrt{2\pi}} \exp\left(-\frac{(k - \mu)^2}{2\sigma^2}\right)$$

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LLT for joint distribution of # vertices and # edges

- Recurrence formulas for # connected graphs [S 70]
- Counting connected graphs inside-out [PW 05]
- Two round exposure and smoothing (for $L_1(d)$) [BC-OK 09]
- Fourier analysis (for joint distribution) [BC-OK 14]

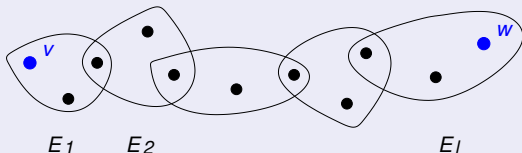
Part II

Random k -uniform Hypergraph $H_k(n, p)$, $k \geq 2$

Standard Notion of Components

Vertex connectivity

- A vertex v is said to be **reachable from a vertex w** if there is a sequence E_1, \dots, E_ℓ of hyperedges such that $v \in E_1$, $w \in E_\ell$ and $|E_i \cap E_{i+1}| \geq 1$ for each $i = 1, \dots, \ell - 1$.



- The reachability is an equivalence relation, and the **equivalence classes** are called **components**

Phase Transition in $H_k(n, p)$

$L_1(d) = \#$ vertices in the largest component, where $d = p \cdot (k-1) \cdot \binom{n-1}{k-1}$

Emergence of giant component

[SCHMIDT-PRUZAN-SHAMIR 85]

- If $d < 1$, whp $L_1(d) = O(\log n)$
- If $d > 1$, whp $L_1(d) = \Theta(n)$

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Local limit theorem for ($\#$ vertices, $\#$ edges) in the giant component

- $(d-1)^3 n \rightarrow \infty$, $(d-1)^3 n = o\left(\frac{\log n}{\log \log n}\right)$ [KAROŃSKI-ŁUCZAK 02]
- $d > 1$ constant [BEHRISCH-COJA-OGHLAN-K. 14]
- $(d-1)^3 n \rightarrow \infty$, $d-1 \rightarrow 0$ [BOLLOBÁS-RIORDAN 14+]

Counting Connected k -uniform Hypergraphs

... with n vertices and m edges

- $m - \frac{n}{k-1} \ll \frac{\log n}{\log \log n}$ [KAROŃSKI-ŁUCZAK 02]
- $m - \frac{n}{k-1} = \Theta(n)$ [BEHRISCH-COJA-OGHLAN-K. 14]
- $m - \frac{n}{k-1} = o(n)$ [BOLLOBÁS-RIORDAN 14+]
- $n^{1/3} \log^2 n \ll m - \frac{n}{2} \ll n$ for $k = 3$ [SATO-WORMALD 14+]

Proof techniques

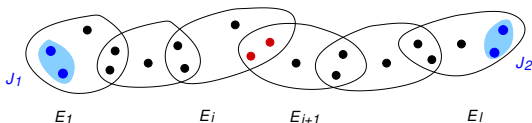
- Combinatorial enumeration [KŁ02]
- Local limit theorem for the giant in $H_k(n, p)$ [BC-OK 14; BR 14+]
- Counting connected graphs inside-out (cores and kernels) [SW 14+]

Higher Order Connectivity

[BOLLOBÁS–RIORDAN 12]

Let $1 \leq j \leq k - 1$.

- A j -element subset J_1 is said to be **reachable** from another j -set J_2 if there is a sequence E_1, \dots, E_ℓ of hyperedges such that $J_1 \subseteq E_1, J_2 \subseteq E_\ell$ and $|E_i \cap E_{i+1}| \geq j$ for each $i = 1, \dots, \ell - 1$.



- The reachability is an equivalence relation on j -sets, and the **equivalence classes** are called j -connected **component**.

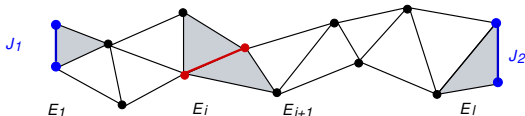
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e.g. $k = 3$, $j = 2$



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Emergence of Giant j -Component

$L_j(d) = \# j$ -sets in the largest j -component, where $d = p \cdot \binom{k}{j} \cdot \binom{n-j}{k-j}$

Theorem

[COOLEY-PERSON-K. 13+]

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Remarks

- Short alternative proof of [Schmidt-Pruzan-Shamir 85]
- Extension of Depth-First Search approach of [Krivelevich-Sudakov 13]
- When $d = 1 + \epsilon$ for $\epsilon \in (0, 1)$,
whp \exists a loose path of length $\Omega(\epsilon^2 n)$

Critical Phase in $H_k(n, p)$

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Proof techniques

- Extension of Breadth-First Search, Galton-Watson branching process and second moment approach of [Bollobás–Riordan 12+]
- Smooth boundary lemma

Part III

Proof Ideas for Supercritical Regime in $H_k(n, p)$

$$k \geq 2, j \geq 1$$

Heuristic for Threshold

- Breadth-First Search process & Galton-Watson branching process

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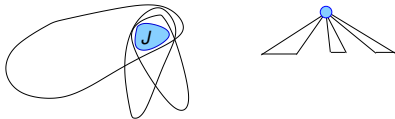


J

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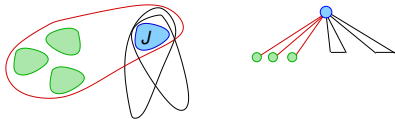
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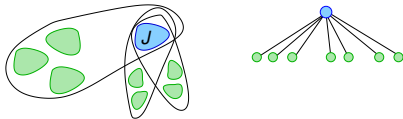
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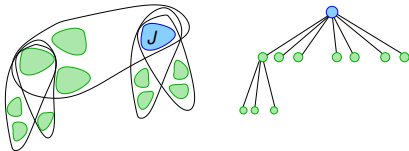


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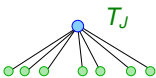
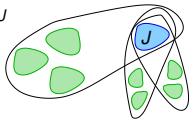
Proof Sketch

(1) Breadth-First Search

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C_J



Given j -set J

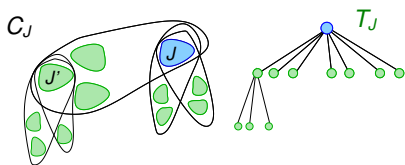
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of j -component C_J

consisting of j -sets as vertices

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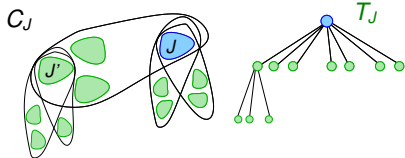
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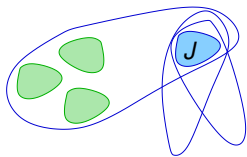
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with offspring distribution $((\binom{k}{j} - 1)Bi(\binom{n-j}{k-j}, p))$

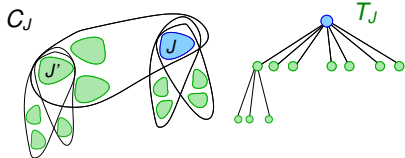


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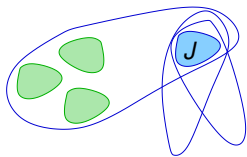
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$$\rightarrow \varrho \sim \frac{2\epsilon}{\binom{k}{j}-1}$$

Proof Sketch – cont.

(3) First moment argument

- Let $N := \# j\text{-sets in 'large' } j\text{-components}$ with $\geq L := \epsilon n^j$ many $j\text{-sets}$
- Using upper and lower couplings with Galton-Watson branching process,

$$\mathbb{E}(N) \sim \frac{2\epsilon}{\binom{k}{j} - 1} \binom{n}{j}$$

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IF we could show

$$\mathbb{E}(N^2) \sim (\mathbb{E}(N))^2,$$

THEN

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(5) Two round exposure

Almost all $j\text{-sets in 'large' } j\text{-components}$ are in a single $j\text{-component}$

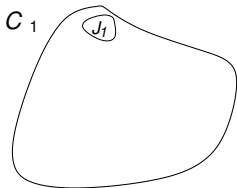
More on Second Moment Argument

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- Fix j -set J_1 and grow its j -component C_1

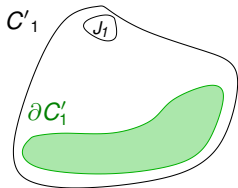


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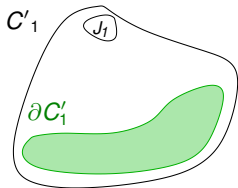
- Fix j -set J_1 and grow its j -component C'_1 until hit **stopping conditions**

$$S_1 = \{ |C'_1| \geq L \text{ or } |\partial C'_1| \geq \epsilon L \}$$



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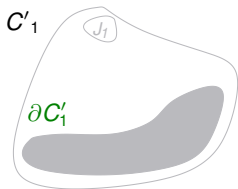
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$$S_1 = \{ |C'_1| \geq L \text{ or } |\partial C'_1| \geq \epsilon L \}$$

$$\text{Then } \mathbb{P}(S_1) \lesssim \frac{2\epsilon}{\binom{k}{j}-1}$$

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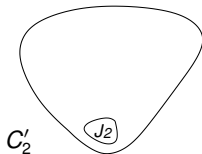


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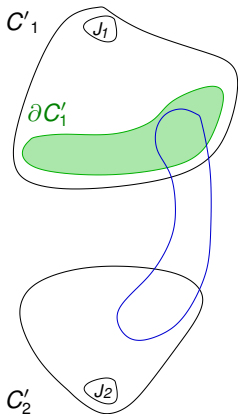
$$\text{Then } \mathbb{P}(S_1) \lesssim \frac{2\epsilon}{\binom{k}{j}-1}$$

- Delete all the vertices in C'_1
& fix a j -set J_2 , grow component C'_2



More on Second Moment Argument

Need to consider # pairs of j -sets in 'large' j -components



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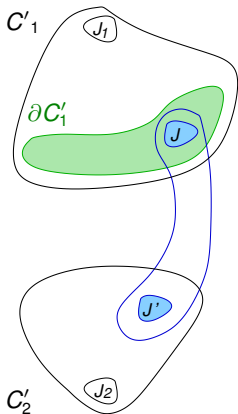
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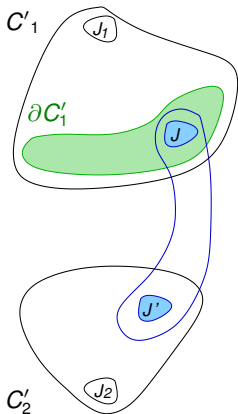
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$$\begin{aligned} \mathbb{P}(e(\partial C'_1, C'_2) \geq 1) \\ \leq p \cdot |\partial C'_1| \cdot |C'_2| \end{aligned}$$

More on Second Moment Argument

Need to consider # pairs of j -sets in 'large' j -components



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$$\text{Then } \mathbb{P}(S_1) \lesssim \frac{2\epsilon}{\binom{k}{j}-1}$$

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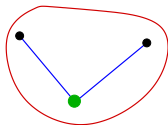
However,

$p \cdot |\partial C'_1| \cdot |C'_2|$ is not the right thing to do

More on Second Moment Argument – cont.

Instead we need

- for $k = 3$, $j = 2$,



$$\mathbb{P}(e(\partial C'_1, C'_2) \geq 1)$$

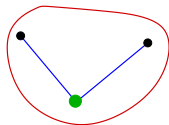
$$\leq \mathbb{E}(\# \text{ 3-sets containing}$$

a pair of 2-sets intersecting at a vertex)

More on Second Moment Argument – cont.

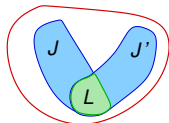
Instead we need

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$$\begin{aligned} & \mathbb{P}(e(\partial C'_1, C'_2) \geq 1) \\ & \leq \mathbb{E}(\# \text{ 3-sets containing} \\ & \quad \text{a pair of 2-sets intersecting at a vertex}) \end{aligned}$$

- for $k \geq 3, j \geq 2$,



$$\begin{aligned} & \mathbb{P}(e(\partial C'_1, C'_2) \geq 1) \\ & \leq \mathbb{E}(\# \text{ k-sets containing} \\ & \quad \text{a pair of j-sets, } J, J', \text{ intersecting at an } \ell\text{-set } L \\ & \quad \text{for some } 0 \leq \ell \leq j - 1) \end{aligned}$$

Boundary Is Smooth

Key lemma

[COOLEY-K.-KOCH 14+]

For every $0 \leq \ell \leq j - 1$, every ℓ -set L ,

$$\# \text{ } j\text{-sets in } \partial C'_j \text{ containing } L \sim \frac{|\partial C'_j|}{\binom{n}{j}} \binom{n-\ell}{j-\ell}$$

'Reasonably Large' Boundary Is Smooth

Key lemma

[COOLEY-K.-KOCH 14+]

Let $\partial C'_1(t)$ denote the collection of j -sets in $\partial C'_1$ after t generations of BFS.

With probability at least $1 - \exp(-\Theta(n^{1/11}))$ the following is true.

For every $0 \leq \ell \leq j - 1$, every ℓ -set L , and every $s_\ell \leq t \leq s_\ell + O(\log n)$,

$$\# \text{ } j\text{-sets in } \partial C'_1(t) \text{ containing } L \sim \frac{|\partial C'_1(t)|}{\binom{n}{j}} \binom{n-\ell}{j-\ell}$$

where $s_\ell := \min\{d : |\partial C'_1(t)| \geq n^{\ell+1/10}\}$.

Open Problems

- (1) What about the number of j -set in the largest j -component at the criticality, i.e. when $d = 1$?
- (2) Is the width of critical window, $(d - 1)^3 n = O(1)$, best possible?
Perhaps $(d - 1)^j n = O(1)$?
- (3) What about the number of j -set in the 2nd largest j -component in the supercritical regime?
- (4) What is the actual distribution of $\#$ j -sets in the largest j -component?
Central limit theorem? Local limit theorem?