

# The duality in the affine action on trees

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## Groups acting on trees:

- ▶ Action of algebraic groups over local fields on their Bruhat-Tits trees
- ▶ Action on rooted trees given by Mealy automata (Burnside groups, groups of intermediate growth, ...)
- ▶ Groups acting on the product of two trees (non-residually finite  $CAT(0)$  groups, finitely presented torsion-free simple groups, ...)

# Automata-transducers

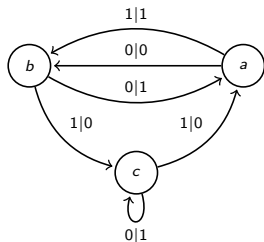
We consider **complete deterministic** Mealy automata.

A finite **automaton** is a quadruple  $A = (S, X, \tau, \pi)$ , where

- ▶  $X$  is a finite set (the input and output alphabet),
- ▶  $S$  is a finite set (the set of states),
- ▶  $\tau : S \times X \rightarrow S$  is the transition map,
- ▶  $\pi : S \times X \rightarrow X$  is the output map.

An automaton can be identified with a **finite directed labeled graph** with vertices  $S$  and arrows

$s \xrightarrow{x|y} t$  iff  $\pi(s, x) = y$  and  $\tau(s, x) = t$ .



# Automaton action on a tree

Let  $X^*$  be the space of finite words over  $X$ . For every state  $s \in S$ , the automaton  $A$  initialized at  $s$  produces the transformation  $A_s : X^* \rightarrow X^*$ :

$$A_s(x_1x_2 \dots x_n) = y_1y_2 \dots y_n$$

for a path  $\textcircled{s} \xrightarrow{x_1|y_1} \textcircled{\phantom{s}} \xrightarrow{x_2|y_2} \textcircled{\phantom{s}} \dots \textcircled{\phantom{s}} \xrightarrow{x_n|y_n} \textcircled{\phantom{s}}$ .

If all  $A_s$  are invertible, they generate an **automaton group**  $G_A$ . A given group is an automaton group iff there is an automaton structure on the group.

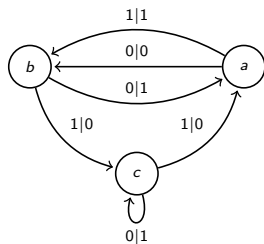
The space  $X^*$  has a natural tree structure. Every automaton group acts by automorphisms on this tree.

# Dual automata

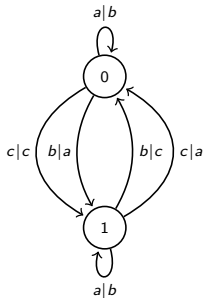
$(S, X, \tau, \pi) \rightsquigarrow$  automaton  $A$  over  $X$  with states  $S$

$(X, S, \tau, \pi) \rightsquigarrow$  **dual automaton**  $\partial A$  over  $S$  with states  $X$

$$s \xrightarrow{x|y} t \text{ in } A \quad \text{iff} \quad x \xrightarrow{s|t} y \text{ in } \partial A$$



$\longleftrightarrow$  **dual**



Automata  $A$  and  $\partial A \rightsquigarrow$  two actions  $S \curvearrowright X^*$  and  $S^* \curvearrowleft X$ .

**Basic question:** What is the relation between these two actions?  
What is the relation between  $G_A$  and  $G_{\partial A}$ ?

**Easy:**  $G_A$  is finite iff  $G_{\partial A}$  is finite

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Connection with Schreier graphs:

The states of  $A^{(n)} = A \circ \dots \circ A$  form a sphere in  $G_A$ .

The  $(\partial A)^{(n)}$  is the **Schreier graph**  $\Gamma_n = \Gamma(G_A, X^n)$ .

# Bireversible automata

There are two standard operations with automata:

taking inverse: if  $s \xrightarrow{x|y} t$  in  $A$ , then  $s^{-1} \xrightarrow{y|x} t^{-1}$  in  $\iota(A)$

taking dual: if  $s \xrightarrow{x|y} t$  in  $A$ , then  $x \xrightarrow{s|t} y$  in  $\partial(A)$

$A$  is invertible iff  $\iota(A)$  is well-defined

$\partial(A)$  is invertible iff  $\iota\partial(A)$  is well-defined

2000 Macedonska, Nekrashevych, Sushchansky: An automaton  $A$  is called **bireversible** if all eight automata

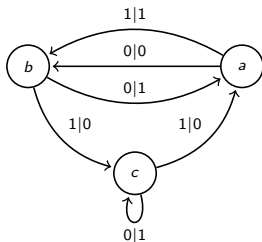
$A, \partial(A), \iota(A), \iota\partial(A), \partial\iota(A), \partial\iota\partial(A), \iota\partial\iota(A), \iota\partial\iota\partial(A) = \partial\iota\partial\iota(A)$

are well-defined (deterministic and complete).

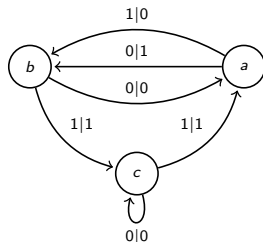


# Aleshin and Bellaterra automata

Aleshin and Bellaterra automata are the only bireversible automata with 3 states over 2 letters generating infinite automaton groups.



Aleshin automaton



Bellaterra automaton

Vorobets<sup>2</sup>:

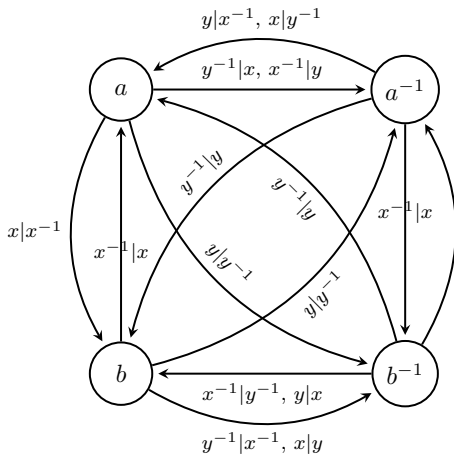
$$G_A = F_3 \text{ free group}$$

$$G_{\partial A} = (\mathbb{Z}_2 \times \mathbb{Z}_2) *_{\mathbb{Z}_2} S_3$$

Muntyan, Nekrashevych, Savchuk:

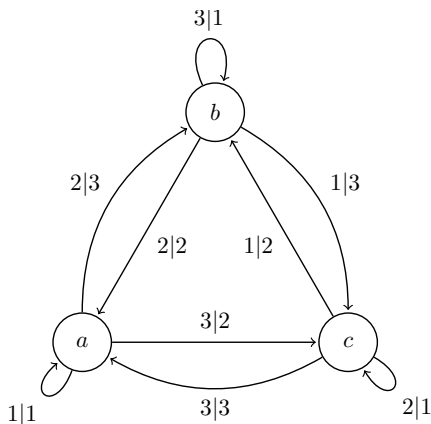
$$G_B = C_2 * C_2 * C_2$$

# One more example



Bondarenko-Kivva:  $A \cong \partial(A) \cong \iota(A)$  and  $G_A = F_2$ .

# Infinitely presented example



Bondarenko-D'Angeli-Rodaro:  $A \cong \partial A \cong \iota(A)$  and  $G_A = \mathbb{Z}_3 \wr \mathbb{Z}$

# Algebraic object behind duality

We can associate a f.p. group with an automaton  $A = (S, X, \tau, \pi)$ :

$$\Gamma_A = \langle S, X \mid sx = yt \text{ for each arrow } s \xrightarrow{x|y} t \text{ in } A \rangle.$$

certain relations in  $\Gamma_A \leftrightarrow$  transitions in  $A$

$$s_1 x_1 x_2 x_3 = y_1 y_2 y_3 s_4 \quad \leftrightarrow \quad s_1 \xrightarrow{x_1|y_1} s_2 \xrightarrow{x_2|y_2} s_3 \xrightarrow{x_3|y_3} s_4$$

$A$  is well-defined:  $SX \subset XS$  in  $\Gamma_A$ .

$\iota(A)$  is well-defined:  $S^{\pm 1}X \subset XS^{\pm 1}$  in  $\Gamma_A$ .

Bireversibility:  $S^{\pm 1}X^{\pm 1} = X^{\pm 1}S^{\pm 1}$  in  $\Gamma_A$ . Therefore, we have  
**group factorization**

$$\Gamma_A = \langle S \rangle \cdot \langle X \rangle$$

# Connection with square complexes

2005 Glasner-Mozes: automata vs square complexes

$$\Delta_A = \left\{ \begin{array}{c} \begin{array}{ccc} & x & \\ \bullet & \xrightarrow{\quad} & \bullet \\ \uparrow & \text{//} & \uparrow \\ s & & t \\ \downarrow & \text{//} & \downarrow \\ \bullet & \xrightarrow{\quad} & \bullet \\ & y & \end{array} & \text{for each arrow } s \xrightarrow{x|y} t \text{ in } A \end{array} \right\}.$$

The complex  $\Delta_A$  is a directed VH square complex with one vertex.

$$\pi_1(\Delta_A) = \langle S, X \mid sx = yt \text{ for each arrow } s \xrightarrow{x|y} t \text{ in } A \rangle.$$

## Proposition (Follows from Wise, Burger-Mozes, Glasner-Mozes)

*Let  $A$  be a finite automaton. The following statements are equivalent:*

- 1  *$A$  is bireversible;*
- 2 *2-cells of  $\Delta_A$  form a 4-way deterministic tileset;*
- 3 *the link of a unique vertex of  $\Delta_A$  is a complete bipartite graph;*
- 4  *$\Delta_A$  is non-positively curved;*
- 5 *the universal cover of  $\Delta_A$  is a direct product of two trees.*

# Properties of $\pi_1(\Delta_A)$

Bireversible property:  $S^{\pm 1}X^{\pm 1} = X^{\pm 1}S^{\pm 1}$  in  $\pi_1(\Delta_A)$ .

## Proposition (Follows from Bridson-Wise)

Let  $A$  be a bireversible automaton. Then:

- 1  $\pi_1(\Delta_A)$  is a torsion-free CAT(0) group.
- 2 The subgroups  $\langle S \rangle$  and  $\langle X \rangle$  of  $\pi_1(\Delta_A)$  are free groups.
- 3  $\pi_1(\Delta_A)$  has exact factorization by  $\langle S \rangle$  and  $\langle X \rangle$ .

Every element  $\gamma \in \pi_1(\Delta_A)$  can be uniquely written in the **normal forms**  $\gamma = gv$  and  $\gamma = uh$  for  $g, h \in \langle S \rangle$  and  $v, u \in \langle X \rangle$ . The correspondence between these two normal forms is given by the automaton  $A^*$ :

$$gv = uh \text{ in } \pi_1(\Delta_A) \quad \text{if and only if} \quad g \xrightarrow{v|u} h \text{ in } A^*.$$

# Automaton groups of bireversible automata

## Proposition (Bondarenko-Kivva)

Let  $A$  be a bireversible automaton. Then

$$G_A \cong F_S/K \quad \text{and} \quad G_{\partial A} \cong F_X/K_{\partial},$$

where  $K$  and  $K_{\partial}$  are the maximal normal subgroups of  $\pi_1(\Delta_A)$  that are contained in  $F_S = \langle S \rangle$  and  $F_X = \langle X \rangle$  respectively.

$G_A, G_{\partial A}$  are finite iff  $\pi_1(\Delta_A)$  is virtually  $F_n \times F_m$  (reducible lattice in the product of two trees)

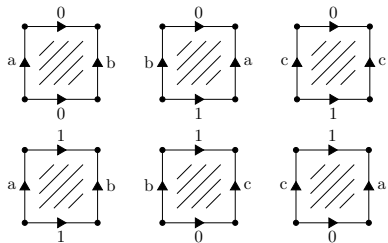
$\pi_1(\Delta_A)$  is just-infinite  $\Rightarrow G_A$  and  $G_{\partial A}$  are free groups

**Question:** What is the relation between  $G_A$  and  $G_{\partial A}$ ?

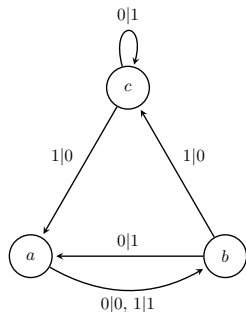
**The transition group:**  $T_A = \pi_1(\Delta_A)/K \times K_{\partial} = G_A \cdot G_{\partial A}$ .



# Aleshin automaton and Wise's example



Wise's complex



Aleshin automaton

2017 Bondarenko-Kivva:  $\pi_1(\Delta_A)$  is non-residually finite

# Affine action on the ring of formal power series

Let  $F$  be a finite field,  $F(t)$  rational function field,  $F[[t]]$  formal power series ring.

The polynomial ring  $F[t]$  has a natural structure of a regular rooted tree with boundary  $F[[t]]$ , given by the sequence of ideals

$$(1) \supset (t) \supset (t^2) \supset (t^3) \supset \dots$$

The affine group

$$\text{Aff}(F[[t]]) = \{\pi_{a,b} : x \mapsto ax + b, a \in F[[t]]^*, b \in F[[t]]\}$$

acts on this tree by automorphisms (self-similar group).

Then:  $\text{Aff}(F[[t]]) \cap \text{FAut}(F[t]) = \text{Aff}(F(t) \cap F[[t]])$

# Example: the lamplighter group

Let  $f(t) = 1 + t \in \mathbb{F}_2[t]$ . Then

$$\pi_{1+t,0}(0 + a_1t + \dots) = 0 + t[(1 + t)(a_1 + a_2t + \dots)]$$

$$\pi_{1+t,0}(1 + a_1t + \dots) = 1 + t[(1 + t)(a_1 + a_2t + \dots) + 1]$$

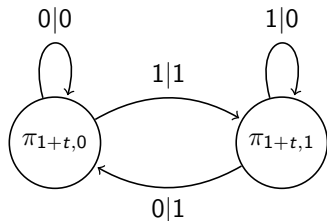
$$\pi_{1+t,1}(0 + a_1t + \dots) = 1 + t[(1 + t)(a_1 + a_2t + \dots)]$$

$$\pi_{1+t,1}(1 + a_1t + \dots) = 0 + t[(1 + t)(a_1 + a_2t + \dots) + 1]$$

$$\begin{aligned} G_{A_f} &= \langle \pi_{1+t,0}, \pi_{1+t,1} \rangle \\ &= \mathbb{F}_2[(1 + t)^{\pm 1}] \rtimes \langle 1 + t \rangle \\ &= \mathbb{Z}_2 \wr \mathbb{Z}, \end{aligned}$$

the lamplighter group

$A_f$  is **not bireversible**



**2005 Silva-Steinberg:** consider  $H[[t]]$  for a finite abelian  $H$ ,  
 $f(t) = 1 - t \in H[[t]]$ ,  $G_{A_f} = H \wr \mathbb{Z}$

**2006 Bartholdi-Sunik:**

- ▶ polynomials over  $\mathbb{Z}/n\mathbb{Z}$ ,  $G_A = (\mathbb{Z}/n\mathbb{Z})^d \wr \mathbb{Z}$
- ▶ the ring  $\mathbb{Z}_n$  of  $n$ -adic integer, multiplication by  $m$ ,  $G_A$  is the Baumslag-Solitar group  $B(1, m)$

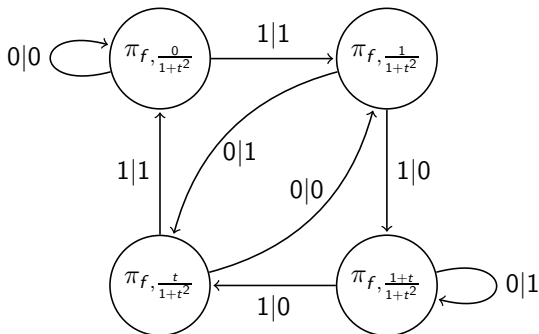
**2006 Bartholdi-Neuhauser-Woess:** consider a finite commutative ring  $R$ ,  $l_i \in R$  satisfying certain conditions, an automaton group

$$\Gamma_d(R) = R[(t + l_1)^{-1}, \dots, (t + l_{d-1})^{-1}, t] \rtimes \langle t + l_1, \dots, t + l_{d-1} \rangle$$

**2017 Skipper-Witzel-Zaremsky:** for every  $n$  a simple group of type  $F_{n-1}$  but not of type  $F_n$

# Example with a rational function

Let  $f(t) = \frac{1+t+t^2}{1+t^2} \in \mathbb{F}_2(t)$ . Then  $A_f$  is bireversible and  $G_{A_f} = (\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2\mathbb{Z}) \wr \mathbb{Z}$  and  $G_{\partial A_f} = \mathbb{Z}/2\mathbb{Z} \wr \mathbb{Z}$ .



# The automaton $A_f$

Let  $f(t) = \frac{p(t)}{q(t)} \in F(t)$  be non-constant. We assume

$$(q(t), t) = 1 \Rightarrow f \in F[[t]] \quad \text{and} \quad \pi_{f,0} \curvearrowright F[[t]]$$

$$(p(t), t) = 1 \Rightarrow f^{-1} \in F[[t]] \quad \text{and} \quad \pi_{f,0} \text{ is invertible}$$

Let  $A_f$  be the automaton for  $\pi_{f,0}$ . The states of  $A_f$  are  $\pi_{f,b}$  for  $b \in B$ , where  $B$  is a finite additive subgroup of  $F(t)$ .

## Proposition

1)  $A_f$  is bireversible iff  $\deg p(t) = \deg q(t)$

2)  $G_{A_f} = B[f, f^{-1}] \rtimes \langle f \rangle \cong (\mathbb{Z}/p\mathbb{Z})^m \wr \mathbb{Z}$  for some  $m \geq 1$

**Question:** What is the dual action?  $\partial A_f \stackrel{?}{=} A_g?$

# Extending the construction

Fix  $g(t) \in F(t)$ . The field extension  $F(t) \supset F(g(t))$  is finite of degree  $d = \deg g(t)$ , here  $1, t, \dots, t^{d-1}$  is a basis. Then every  $f(t) \in F(t)$  can be written as

$$\begin{aligned} f(t) &= r_0(g(t)) + t \cdot r_1(g(t)) + \dots + t^{d-1} r_{d-1}(g(t)) = \\ &= \sum_{n \geq -k} a_n(t) g^n(t), \text{ where } a_n(t) \in F_{<d}[t]. \end{aligned}$$

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$R_g = \{f(t) \in F(t) : f(t) = \sum_{n \geq 0} a_n(t) g^n(t)\}$  is a subring of  $F(t)$ .

We can consider **the action of  $\text{Aff}(R_g)$  on the tree  $T_g$**  given by the ideals  $(g^n(t)), n \geq 0$ .



Let  $f(t) = \frac{p_1(t)}{q_1(t)}, g(t) = \frac{p_2(t)}{q_2(t)} \in F(t)$  be such that  $f(t) \in R_g$ . We can consider the action of  $f(t)$  on the tree  $T_g$  by multiplication. Let  $A_{(f,g)}$  be the corresponding automaton.

If  $g(t) = t$  then  $A_{(f,g)} = A_f$  is just standard automaton for  $f(t)$ .

## Theorem

- 1)  $A_{(f,g)}$  is bireversible if and only if  $p_1, p_2, q_1, q_2$  are pairwise coprime and either  $\deg p_1(t) = \deg q_1(t)$  or  $\deg p_2(t) = \deg q_2(t)$ .
- 2)  $G_{A_{(f,g)}}$  is a lamplighter group.

# Automata $A_{(f,g)}$

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**Inverse & Dual:**  $\iota(A_{(f,g)}) = A_{(f^{-1},g)}$  and  $\partial(A_{(f,g)}) = A_{(g^{-1},f^{-1})}$

Eight automata correspond to the pairs  $(f^{\pm 1}, g^{\pm 1})$  and  $(g^{\pm 1}, f^{\pm 1})$ .

# Transition group of $A_{(f,g)}$

Let  $f(t), g(t) \in F(t)$  and consider two subsets of  $\text{Aff}(F(t))$ :

$$X_f = \{\pi_{f,a}, a \in S_f\} \quad \text{and} \quad X_g = \{\pi_{g,b}, b \in S_g\},$$

where  $S_{\frac{p(t)}{q(t)}} = \left\{ a(t)/q(t) : a(t) \in F[t], \deg a(t) < \deg \frac{p(t)}{q(t)} \right\}$ .

The automaton  $A_{(f,g)}$  has the states  $X_f$  over the alphabet  $X_g$ ,

$$\pi_{f,a} \xrightarrow{\pi_{g,b} | \pi_{g,b'}} \pi_{f,a'} \quad \text{iff} \quad \pi_{f,a} \circ \pi_{g,b} = \pi_{g,b'} \circ \pi_{f,a'}.$$

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Then

$$\begin{aligned} G_A &= \langle X_f \rangle = S_f[f, f^{-1}] \rtimes \langle f \rangle \\ G_{\partial A} &= \langle X_g \rangle = S_g[g, g^{-1}] \rtimes \langle g \rangle \\ T(A_{(f,g)}) &= \langle X_f, X_g \rangle = \langle X_f \rangle \langle X_g \rangle = G_A \cdot G_{\partial A} \\ &= (S_f[f, f^{-1}] + S_g[g, g^{-1}]) \rtimes \langle f, g \rangle \end{aligned}$$

# The rings of $S$ -integers

Let  $S$  be a finite set of places of  $F(t)$  and  $\mathcal{O}_S$  the ring of  $S$ -integers of  $F(t)$ . We can find  $g(t) \in F(t)$  such that

$$\mathcal{O}_S \cap S_g[g, g^{-1}] = \{0\} \quad \Rightarrow \quad \text{Aff}(\mathcal{O}_S) \cap \text{Aff}(S_g[g, g^{-1}]) = \{id\}$$

## Theorem

*The group  $\text{Aff}(\mathcal{O}_S)$  can be realized by a bireversible automaton*

**Corollary:** Bireversible automata  $A$  such that  $G_A$  is finitely presented, while  $G_{\partial A}$  is not.

**Corollary:** A basic construction of

- ▶ a family of bireversible automata
- ▶ irreducible lattices in the product of two trees
- ▶ non-residually finite CAT(0) groups
- ▶ directed VH square complexes

**Question:** What is the nature of  $\pi_1(\Delta_{A_{(f,g)}})$ ?

**Corollary:** A basic construction of

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**Question:** What is the nature of  $\pi_1(\Delta_{A_{(f,g)}})$ ?

**Conjecture:** Groups generated by bireversible automata are linear

Thank You for attention!