

### Exercise sheet 3

Exercises for the exercise session on 3 November 2020

**Problem 3.1.** Let  $A, B, C$  be events in a probability space such that

- $A$  and  $B$  are independent;
- $B$  and  $C$  are independent;
- $A \cap B$  and  $C$  are independent.

Prove that  $A$  and  $B \cap C$  are independent as well. Deduce that any *minimal* dependency graph  $D$  for events  $A_1, \dots, A_n$  has the following properties.

- Every vertex  $i$  of  $D$  is either isolated (i.e. there are no edges involving  $i$ ) or has at least one outgoing edge (i.e. an edge  $(i, j)$ ).
- Either  $D$  is edgeless or contains a directed cycle (this cycle might consist of only two edges  $(i, j)$  and  $(j, i)$ ).

Use these properties to show that in the following setup, Lovász Local Lemma *cannot* be applied (regardless of the choice of the dependency graph  $D$ ). Let  $\Omega$  be the set of all sequences  $s_1, \dots, s_n$  with  $s_i \in \{1, \dots, 6\}$  such that  $s_1 + \dots + s_n$  is even. Choose an element from  $\Omega$  uniformly at random (i.e. we are considering the outcome of throwing a fair die  $n$  times, conditioned on the event that the sum of values is even). Denote by  $A_i$  the event that  $s_i$  is odd.

**Problem 3.2.** Let  $A_1, \dots, A_n$  be events in a probability space and let  $D = ([n], E)$  be a directed graph. We say that  $D$  is a *negative dependency graph* for  $A_1, \dots, A_n$  if the following holds.

For every  $J \subseteq [n]$  and  $i \in [n] \setminus J$ , if  $(i, j) \notin E$  for all  $j \in J$  and  $\mathbb{P} \left[ \bigwedge_{j \in J} \overline{A_j} \right] > 0$ , then  $\mathbb{P} \left[ A_i \mid \bigwedge_{j \in J} \overline{A_j} \right] \leq \mathbb{P}(A_i)$ .

Replacing ‘dependency graph’ by ‘negative dependency graph’ in the statement of the Lovász Local Lemma results in the so-called *Lopsided Lovász Local Lemma*.

- Sketch how the proof of the Lovász Local Lemma from the lecture generalises to a proof of the Lopsided Lovász Local Lemma.
- Consider the following setup. Let  $\Omega$  be the set of all permutations of  $[n]$ . We choose a permutation  $\sigma$  uniformly at random and denote, for each  $i \in [n]$ , by  $A_i$  the event that  $\sigma(i) = i$ .

Follow the arguments sketched below to prove that the edgeless graph on  $[n]$  is a negative dependency graph.

It suffices to show (why?) that

$$\mathbb{P} \left[ \bigwedge_{j \in J} \overline{A_j} \mid A_i \right] \leq \mathbb{P} \left[ \bigwedge_{j \in J} \overline{A_j} \right]$$

for all  $J \subset [n]$  and  $i \in [n] \setminus J$  with  $\mathbb{P}[A_i] > 0$ . Show that this is equivalent to

$$n \left| A_i \wedge \bigwedge_{j \in J} \overline{A_j} \right| \leq \left| \bigwedge_{j \in J} \overline{A_j} \right|.$$

To prove this inequality, for each permutation  $\sigma \in A_i \wedge \bigwedge_{j \in J} \overline{A_j}$ , define  $\sigma_1, \dots, \sigma_n \in \bigwedge_{j \in J} \overline{A_j}$  and prove that  $\sigma_k \neq \tau_l$  as soon as  $k \neq l$  or  $\sigma \neq \tau$ .

**Problem 3.3.** We say that a hypergraph  $H = (V, E)$  is *2-colourable* if there exists a colouring of  $V$  by two colours so that no edge in  $E$  is monochromatic.

Let  $k \geq 2$  be given.

- (a) Let  $H$  be a hypergraph in which every edge has at least  $k$  vertices. Suppose that each edge of  $H$  intersects at most  $d \geq 1$  other edges. Prove that  $H$  is 2-colourable if  $e(d+1)2^{1-k} \leq 1$ .
- (b) Suppose that  $H$  is a hypergraph in which each edge has at least  $k$  vertices. For each edge  $f$  and each  $j \geq k$ , denote by  $d_{f,j}$  the number of edges of size  $j$  that intersect  $f$ . Prove that if for each edge  $f$  of  $H$

$$8 \sum_{j \geq k} \frac{d_{f,j}}{2^j} \leq 1,$$

then  $H$  is 2-colourable.

**Problem 3.4.** Let  $G$  be a graph and let  $d \geq 1$ . Suppose that for every vertex  $v$ , there exists a list  $S(v)$  of precisely  $\lceil 2ed \rceil$  ‘admissible’ colours such that no colour in  $S(v)$  is admissible for more than  $d$  neighbours of  $v$ . Prove that there is a ‘proper’ colouring of  $G$  (i.e. no two adjacent vertices have the same colour) assigning to each vertex an admissible colour.

*Hint.* The fewer vertices and colours play a role in the probability of a ‘bad’ event  $A$ , the simpler the expression for  $\mathbb{P}[A]$  will be.

**Problem 3.5.** Suppose we place  $n$  balls in  $n$  bins, where each ball chooses its bin uniformly at random and independently from the other balls.

- (a) Prove that for each  $\varepsilon > 0$ ,

$$\mathbb{P} \left[ \exists \text{ a bin with at least } \left( \frac{2}{3} + \varepsilon \right) \ln n \text{ balls} \right] = o(1).$$

- (b) By how much can we decrease the value  $\left( \frac{2}{3} + \varepsilon \right) \ln n$  in (a) so that we can still prove (by the same type of arguments as in (a)) that the probability is  $o(1)$ ?
- (c) If we have  $n^2$  balls in total, for what  $k = k(n)$  can we prove that

$$\mathbb{P}[\exists \text{ a bin with at most } k \text{ balls}] = o(1)?$$

**Problem 3.6.** Define the set  $S \subset \mathbb{N}$  by letting each number  $n$  be in  $S$  with probability  $1/2$  independently.

(a) For  $k, l \in \mathbb{N}$ , we set

$$w_l(k) = \left\lceil \frac{\ln(kl2^{k-1})}{\ln 2} \right\rceil.$$

Denote by  $A_l$  the event that there is a  $k \geq 2$  such that  $S$  contains an arithmetic progression of the form

$$k - b, k, k + b, \dots, k + (w_l(k) - 2)b.$$

Prove that  $\mathbb{P}[A_l] \leq 1/l$  and deduce from this that with probability 1,  $S$  does not contain an arithmetic progression of infinite length.

(b) Prove that

$$\mathbb{P} \left[ \lim_{n \rightarrow \infty} \frac{|S \cap [n]|}{n} = \frac{1}{2} \right] = 1.$$

To that end, for fixed  $\varepsilon > 0$  and  $n$ , use the Chernoff bounds to find an upper bound for

$$\mathbb{P} \left[ \left| \frac{|S \cap [n]|}{n} - \frac{1}{2} \right| \geq \varepsilon \right]$$

and apply a union bound to show that

$$\mathbb{P} \left[ \exists n \geq n_0 \text{ with } \left| \frac{|S \cap [n]|}{n} - \frac{1}{2} \right| \geq \varepsilon \right] \stackrel{n_0 \rightarrow \infty}{\equiv} o(1).$$

Where does this strategy fail when we use Chebyshev's inequality instead of Chernoff bounds?

(*Note.* Szemerédi's Theorem states that each  $A \subset \mathbb{N}$  with

$$\limsup_{n \rightarrow \infty} \frac{|A \cap [n]|}{n} > 0$$

contains infinitely many arithmetic progressions of length  $k$  for *every*  $k$ . Thus, Problem 3.6 shows that with probability 1,  $S$  contains arbitrarily long arithmetic progressions, but no arithmetic progression of infinite length.)