## Advanced and algorithmic graph theory



Summer term 2022

## Exercise sheet 1

Exercises for the exercise session on 21/03/2022

**Problem 1.1.** Let G be a connected graph. Prove that the following statements are equivalent for an edge  $e \in E(G)$ .

- (i) G e is not connected;
- (ii) No cycle in G contains e;
- (iii) Every spanning tree of G contains e;
- (iv) Every spanning tree of G constructed by depth-first search contains e (independently of the order in which we check vertices in the FOR-loops).

**Problem 1.2.** Let G be a connected weighted graph. For i = 1, 2, denote by  $T_i$  the graph that is generated by Algorithm i from the lecture. Reminder:

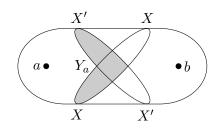
- Alg. 1: Start with  $T_1 = (V(G), \emptyset)$ . Recursively add to  $T_1$  edges of smallest weight that do not create a cycle.
- Alg. 2: Start with  $T_2 = G$  and recursively delete edges of largest weight that do not disconnect  $T_2$ .

Prove that both  $T_1$  and  $T_2$  are spanning trees of G and that both have smallest total weight among all spanning trees of G.

**Problem 1.3.** Let G be a graph and let a, b be two distinct vertices of G. Suppose that each of the vertex sets  $X, X' \subseteq V(G) \setminus \{a, b\}$  is an *a-b* separator. Denote by  $C_a$  and  $C_b$  the component of G - X that contains a and b, respectively. Define  $C'_a$  and  $C'_b$  analogously for X'. Prove that the sets

$$Y_a := (X \cap C'_a) \cup (X \cap X') \cup (X' \cap C_a),$$
  
$$Y_b := (X \cap C'_b) \cup (X \cap X') \cup (X' \cap C_b)$$

are a-b separators as well.



**Problem 1.4.** Let  $G, a, b, X, X', Y_a, Y_b$  be as in Problem 1.3.

- (a) Prove that X is a minimal a-b separator (with respect to containment, i.e. no proper subset of X is an a-b separator) if and only if each vertex in X has neighbours in both  $C_a$  and  $C_b$ .
- (b) Suppose that both X and X' have smallest size among all a-b separators in  $V \setminus \{a, b\}$ . Prove that  $Y_a$  and  $Y_b$  are then also smallest a-b separators in  $V \setminus \{a, b\}$ .
- (c) Give an example for which X and X' are minimal a-b separators (w.r.t. containment), but  $Y_a$  and  $Y_b$  are not minimal.

**Problem 1.5.** Let G be k-connected, where  $k \ge 2$ . Prove that G contains a cycle of length at least min $\{2k, |G|\}$ . Show that this statement is best possible in the sense that for every  $k \ge 2$ , there is a k-connected graph with no cycle of length at least min $\{2k + 1, |G|\}$ .

**Problem 1.6.** Find all mistakes in the following 'proof' of the set version of Menger's theorem. What statement(s) would be necessary to prove in order to complete the proof?

Let S be a smallest A-B separator. We say that a component C of G-S meets A if C contains a vertex of A. Denote by  $G_A$  the graph that G induces on the union of S and all vertex sets of components of G-S that meet A. Define  $G_B$  analogously. By the choice of S and induction,  $G_A$  contains |S| disjoint A-S paths, while  $G_B$ contains |S| disjoint S-B paths. Joining these paths yields the desired set of |S|disjoint A-B paths.