# $\mathcal{H}^{2}$-matrices with adaptive bases 

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Goal: Treat certain large dense matrices in linear complexity.
Model problem: Discretization of integral operators of the form

$$
\mathcal{K}[u](x)=\int_{\Omega} \kappa(x, y) u(y) d y
$$

Discretization: Galerkin method with FE basis $\left(\varphi_{i}\right)_{i=1}^{n}$ leads to a $\operatorname{matrix} K \in \mathbb{R}^{n \times n}$ with

$$
K_{i j}=\left\langle\varphi_{i}, \mathcal{K}\left[\varphi_{j}\right]\right\rangle_{L^{2}}=\int_{\Omega} \int_{\Omega} \varphi_{i}(x) \kappa(x, y) \varphi_{j}(y) d y d x . \quad(1 \leq i, j \leq n)
$$

Problem: $K$ is dense.
Idea: Use properties of the kernel function to approximate $K$.
Examples: Wavelets, multipole expansion, interpolation.

Idea: We consider a subdomain $\tau \times \sigma \subseteq \Omega \times \Omega$ and replace the kernel function $\kappa$ by an interpolant

$$
\tilde{\kappa}^{\tau, \sigma}(x, y):=\sum_{\nu=1}^{k} \sum_{\mu=1}^{k} \kappa\left(x_{\nu}^{\tau}, x_{\mu}^{\sigma}\right) \mathcal{L}_{\nu}^{\tau}(x) \mathcal{L}_{\mu}^{\sigma}(y) . \quad(k \ll n)
$$

Result: Local matrix in factorized form

$$
\begin{aligned}
K_{i j}^{\tau, \sigma} & :=\int_{\tau} \int_{\sigma} \varphi_{i}(x) \kappa(x, y) \varphi_{j}(y) d y d x \approx \int_{\tau} \int_{\sigma} \varphi_{i}(x) \tilde{\kappa}^{\tau, \sigma}(x, y) \varphi_{j}(y) d y d x \\
& =\sum_{\nu=1}^{k} \sum_{\mu=1}^{k} \underbrace{\kappa\left(x_{\nu}^{\tau}, x_{\mu}^{\sigma}\right)}_{=: S_{\nu \mu}^{\tau, \sigma}} \underbrace{\int_{\tau}^{\varphi_{i}} \varphi_{i}(x) \mathcal{L}_{\nu}^{\tau}(x) d x}_{=: V_{i \nu}^{\tau}} \underbrace{\int_{\sigma} \varphi_{j}(y) \mathcal{L}_{\mu}^{\sigma}(y) d y}_{=: V_{j \mu}^{\sigma}} \\
& =\left(V^{\tau} S^{\tau, \sigma}\left(V^{\sigma}\right)^{\top}\right)_{i j}
\end{aligned}
$$

Problem: Typical kernel functions are not globally smooth.
Idea: They are locally smooth outside of the diagonal $x=y$.
Examples: $1 /\|x-y\|, \log \|x-y\|$.
Error bound: For $m$-th order tensor interpolation on axis-parallel boxes $B^{\tau} \supseteq \tau$ and $B^{\sigma} \supseteq \sigma$, we get

$$
\left\|\tilde{\kappa}^{\tau, \sigma}-\kappa\right\|_{L^{\infty}\left(B^{\tau} \times B^{\sigma}\right)} \leq C_{\mathrm{apx}}(m)\left(1+\frac{2 \operatorname{dist}\left(B^{\tau}, B^{\sigma}\right)}{\operatorname{diam}\left(B^{\tau} \times B^{\sigma}\right)}\right)^{-m}
$$

Rank: $m$ interpolation points per coordinate $\Rightarrow k=m^{d}$.
Admissibility: Uniform exponential convergence requires

$$
\operatorname{diam}\left(B^{\tau} \times B^{\sigma}\right) \leq 2 \eta \operatorname{dist}\left(B^{\tau}, B^{\sigma}\right)
$$



Approach: Split $\Omega \times \Omega$ into admissible subdomains and a small remainder.

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Cluster tree: Hierarchy of subdomains of $\Omega$
Nearfield: Small blocks, stored in standard format.
Farfield: Admissible blocks, stored in $V^{\tau} S^{\tau, \sigma}\left(V^{\sigma}\right)^{\top}$ format. Partition: Collection of farand nearfield blocks.

Construction of cluster tree and block partition can be accomplished by general algorithms in $\mathcal{O}(n)$ operations.

Goal: Treat the cluster basis $\left(V^{\tau}\right)_{\tau \in \mathcal{T}}$ efficiently.
Idea: For a cluster $\tau$ and $\tau^{\prime} \in \operatorname{sons}(\tau)$ we have

$$
\mathcal{L}_{\nu}^{\tau}=\sum_{\nu^{\prime}=1}^{k} \mathcal{L}_{\nu}^{\tau}\left(x_{\nu^{\prime}}^{\tau^{\prime}}\right) \mathcal{L}_{\nu^{\prime}}^{\tau^{\prime}} . \quad(1 \leq \nu \leq k)
$$

Nested bases: Consider $i \in \tau^{\prime} \subseteq \tau$. Setting $T_{\nu^{\prime} \nu}^{\tau^{\prime}}:=\mathcal{L}_{\nu}^{\tau}\left(x_{\nu^{\prime}}^{\tau^{\prime}}\right)$, we find

$$
V_{i \nu}^{\tau}=\int_{\tau} \varphi_{i}(x) \mathcal{L}_{\nu}^{\tau}(x) d x=\sum_{\nu^{\prime}=1}^{k} T_{\nu^{\prime} \nu}^{\tau^{\prime}} \int_{\tau} \varphi_{i}(x) \mathcal{L}_{\nu^{\prime}}^{\tau^{\prime}}(x) d x=\left(V^{\tau^{\prime}} T^{\tau^{\prime}}\right)_{i \nu}
$$

Consequence: Store $V^{\tau}$ only for leaves and use $T^{\tau^{\prime}}$ for all other clusters.
Complexity: Storage $\mathcal{O}(n k)$, matrixvector $\mathcal{O}(n k)$.


Problem: Efficiency depends on the rank $k$. Can we reduce this rank by eliminating redundant expansion functions?

Idea: Orthogonalize columns of $V^{\tau}$.
Gram-Schmidt: Find $\tilde{V}^{\tau}=V^{\tau} Z^{\tau}$ such that $\tilde{V}^{\tau}=V^{\tau}$ $\left(\tilde{V}^{\tau}\right)^{\top} \tilde{V}^{\tau}=I$ holds.

Nested structure: Using $P^{\tau}$ with $V^{\tau}=\tilde{V}^{\tau} P^{\tau}$, we find

$$
V^{\tau}=\binom{V^{\tau_{1}} T^{\tau_{1}}}{V^{\tau_{2}} T^{\tau_{2}}}=\binom{\tilde{V}^{\tau_{1}} P^{\tau_{1}} T^{\tau_{1}}}{\tilde{V}^{\tau_{2}} P^{\tau_{2}} T^{\tau_{2}}}=\left(\begin{array}{ll}
\tilde{V}^{\tau_{1}} & \\
& \tilde{V}^{\tau_{2}}
\end{array}\right)\binom{P^{\tau_{1}} T^{\tau_{1}}}{P^{\tau_{2}} T^{\tau_{2}}}
$$

so it is sufficient to orthogonalize a $(2 k) \times k$ matrix. The resulting cluster basis is nested due to

$$
\tilde{V}^{\tau}=V^{\tau} Z^{\tau}=\binom{\tilde{V}^{\tau_{1}}\left(P^{\tau_{1}} T^{\tau_{1}} Z^{\tau}\right)}{\tilde{V}^{\tau_{2}}\left(P^{\tau_{2}} T^{\tau_{2}} Z^{\tau}\right)}=\binom{\tilde{V}^{\tau_{1}} \tilde{T}^{\tau_{1}}}{\tilde{V}^{\tau_{2}} \tilde{T}^{\tau_{2}}}
$$

Result: Complexity $\mathcal{O}\left(n k^{2}\right)$.

Example: Unit sphere, piecewise constant trial and test functions.

|  | Interpolation |  |  | Orthogonalized |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $n$ | Build | MVM | Mem/n | Expl | Impl | MVM | Mem/n |
| 2048 | 6 | 0.14 | 49.2 | 3 | 8 | 0.07 | 17.0 |
| 8192 | 26 | 0.54 | 45.6 | 15 | 35 | 0.31 | 21.1 |
| 32768 | 126 | 2.45 | 49.6 | 107 | 149 | 1.37 | 22.0 |
| 131072 | 765 | 10.33 | 53.4 | 469 | 774 | 5.48 | 22.0 |
| 524288 | 3347 | 43.08 | 53.8 | 2045 | 3473 | 23.81 | 22.0 |

$\mathcal{H}^{2}$-matrix constructed by cubic interpolation and subsequent orthogonalization, relative error $\leq 10^{-3}$.

Hardware: SunFire 6800 with 900 MHz UltraSparc IIIcu processors
Software: HLIB, see http://www.hlib.org

Problem: Interpolation does not take into account that a subset of polynomials may be sufficient to approximate the kernel function.

Idea: Use local singular value decompositions to find optimal cluster bases $\Rightarrow$ Take kernel and geometry into account.

Important: Since the cluster bases are nested, the father clusters depend on the son clusters.


Consequence: If $\tau \times \sigma$ is a farfield block and $\tau^{\prime} \in \operatorname{sons}(\tau)$, we have

$$
\left.\tilde{K}\right|_{\tau^{\prime} \times \sigma}=\left.\left(V^{\tau} S^{\tau, \sigma}\left(V^{\sigma}\right)^{\top}\right)\right|_{\tau^{\prime} \times \sigma}=V^{\tau^{\prime}} T^{\tau^{\prime}} S^{\tau, \sigma}\left(V^{\sigma}\right)^{\top},
$$

i.e., the choice of $V^{\tau^{\prime}}$ influences all farfield blocks connected to $\tau$ or its other ancestors.


Block row: $R^{\tau}:=\left\{\sigma: \exists \tau^{+} \supseteq \tau: \tau^{+} \times \sigma\right.$ is admissible $\}$.
Goal: Find an orthogonal matrix $V^{\tau} \in \mathbb{R}^{\tau \times k^{\tau}}$ that maximizes

$$
\sum_{\sigma \in R^{\tau}}\left\|\left.\left(V^{\tau}\right)^{\top} K\right|_{\tau \times \sigma}\right\|_{F}^{2}
$$

Optimal solution: Use eigenvectors of corresponding Gram matrix.
Nested structure: Ensured by bottom-up recursion and projections.
Result: Algorithm with complexity $\mathcal{O}\left(n k^{2}\right)$, truncation error can be controlled directly.

Example: Unit sphere, piecewise constant trial and test functions.

|  | Interpolation |  |  | Recompressed |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $n$ | Build | MVM | Mem/n | Build | MVM | Mem/n |
| 2048 | 6 | 0.14 | 49.2 | 11 | 0.01 | 3.7 |
| 8192 | 26 | 0.54 | 45.6 | 51 | 0.11 | 4.4 |
| 32768 | 128 | 2.41 | 49.6 | 210 | 0.45 | 4.6 |
| 131072 | 762 | 10.34 | 53.4 | 883 | 2.07 | 4.7 |
| 524288 | 3371 | 43.37 | 53.8 | 3818 | 9.77 | 5.2 |

$\mathcal{H}^{2}$-matrix constructed by cubic interpolation.
Recompression with relative error bound $10^{-3}$, maximal rank 16 .
Hardware: SunFire 6800 with 900 MHz UltraSparc IIIcu processors.
Software: HLIB, see http://www.hlib.org

Example: Grid "Crank shaft" of the Netgen package.


| $n$ | 1768 | 10126 | 113152 |
| ---: | ---: | ---: | ---: |
| Build | 20 | 200 | 2255 |
| Mem/n | 5.7 | 7.6 | 5.6 |
| Near/n | 4.2 | 4.7 | 3.4 |
| MVM | 0.02 | 0.24 | 2.18 |
| Rel. error | $0.05 \%$ | $0.07 \%$ | $0.08 \%$ |
| DMem/n | 13.8 | 79.1 | 884.0 |

$$
\begin{aligned}
b(v, \phi)= & \int_{\Gamma} \int_{\Gamma}\langle\nabla \times \phi(y), v(x)\rangle\left\langle\nabla_{x} \Phi(x, y), n(x)\right\rangle d y d x \\
& -\int_{\Gamma} \int_{\Gamma}\langle\nabla \times \phi(y), n(x)\rangle\left\langle\nabla_{x} \Phi(x, y), v(x)\right\rangle d y d x
\end{aligned}
$$

|  |  |  |  |
| ---: | ---: | ---: | ---: |
|  | $n$ | 42432 | 169728 |
| Direct | 1554 | 5178 | 37632 |
| MVM | 8.3 | 32.5 | 1428 |
| Mem/DoF | 168 | 165 | 8.0 |
| Recompressed | 1700 | 6090 | 183 |
| MVM | 0.8 | 2.6 | 0.7 |
| Mem/DoF | 8.0 | 6.5 | 8.0 |
| Rel. error | $8 \times 10^{-3}$ |  | $1 \times 10^{-3}$ |

$\mathcal{H}^{2}$-matrices: Nested structure leads to complexity $\mathcal{O}(n k)$ for storage requirements and matrix-vector multiplication.

Black box: The approximation is constructed by using only kernel evaluations $\kappa\left(x_{\nu}^{\tau}, x_{\mu}^{\sigma}\right)$ and works for all asymptotically smooth kernel functions.

Recompression: Use $\mathcal{O}\left(n k^{2}\right)$ algorithm to reduce the rank $k$, find a quasi-optimal cluster basis adapted to geometry, discretization and kernel function.

Work in progress: $\mathcal{H}^{2}$-matrix arithmetics, HCA.

Goal: Perform matrix addition and multiplication efficiently.
Addition and scaling: For an admissible block $\tau \times \sigma$, we have

$$
\left.(A+\lambda B)\right|_{\tau \times \sigma}=V^{\tau}\left(S_{A}^{\tau, \sigma}+\lambda S_{B}^{\tau, \sigma}\right)\left(V^{\sigma}\right)^{\top}
$$

so $\mathcal{H}^{2}$-matrices form a vector space. Addition and scaling can be performed in $\mathcal{O}(n k)$ operations.


Multiplication: Defined recursively.

$$
\left(\begin{array}{ll}
A_{11} & A_{12} \\
A_{21} & A_{22}
\end{array}\right)\left(\begin{array}{ll}
B_{11} & B_{12} \\
B_{21} & B_{22}
\end{array}\right)=\left(\begin{array}{ll}
A_{11} B_{11}+A_{12} B_{21} & A_{11} B_{12}+A_{12} B_{22} \\
A_{21} B_{11}+A_{22} B_{21} & A_{21} B_{12}+A_{22} B_{22}
\end{array}\right)
$$

Use truncation if product does not match prescribed format.


Situation: Let $A=V^{\tau} S_{A}^{\tau, \varrho}\left(V^{\varrho}\right)^{\top}, B$ not admissible.
Compute best approximation of $A B$ in the form $C=V^{\tau} S_{C}^{\tau, \sigma}\left(V^{\sigma}\right)^{\top}$.
Orthogonality: If $V^{\tau}$ and $V^{\sigma}$ are orthogonal:

$$
S_{C}^{\tau, \sigma}=\left(V^{\tau}\right)^{\top} A B V^{\sigma}=\left(V^{\tau}\right)^{\top} V^{\tau} S_{A}^{\tau, \varrho}\left(V^{\varrho}\right)^{\top} B V^{\sigma}=S_{A}^{\tau, \varrho} \underbrace{\left(V^{\varrho}\right)^{\top} B V^{\sigma}}_{=: \tilde{S}_{B}^{\varrho, \sigma}} .
$$

Problem: Computing $\left(V^{\varrho}\right)^{\top} B V^{\sigma}$ directly is too expensive.
Idea: Prepare $\tilde{S}_{B}^{\varrho, \sigma}:=\left(V^{\varrho}\right)^{\top} B V^{\sigma}$ in advance.
Matrix forward transformation: Since the cluster bases are nested, all matrices $\tilde{S}_{B}^{\varrho, \sigma}$ can be prepared in $\mathcal{O}\left(n k^{3}\right)$ operations.


Situation: Let $A=V^{\tau} S_{A}^{\tau, Q}\left(V^{\varrho}\right)^{\top}$ and $B=V^{\varrho} S_{B}^{o, \sigma}\left(V^{\sigma}\right)^{\top}$.
Advantage: The product $A B$ already has the factorized form

$$
A B=V^{\tau} S_{A}^{\tau, Q}\left(V^{\varrho}\right)^{\top} V^{\varrho} S_{B}^{\varrho, \sigma}\left(V^{\sigma}\right)^{\top}=V^{\tau}\left(S_{A}^{\tau, Q} S_{B}^{\varrho, \sigma}\right)\left(V^{\sigma}\right)^{\top} .
$$

Problem: Converting $A B$ directly to the format of $C$ will lead to inacceptable complexity.
Idea: Store result in $\tilde{S}_{C}^{\tau, \sigma}:=S_{A}^{\tau, Q} S_{B}^{\rho, \sigma}$ and distribute after the multiplication is complete.

Matrix backward transformation: Since the cluster bases are nested, all matrices $\tilde{S}_{C}^{\tau, \sigma}$ can be distributed in $\mathcal{O}\left(n k^{3}\right)$ operations.


Situation: Let $B=V^{Q} S_{B}^{\rho, \sigma}\left(V^{\sigma}\right)^{\top}$.
Approach: Split $B$ and proceed by recursion:

$$
\left.B\right|_{\varrho^{\prime} \times \sigma^{\prime}}=V^{\varrho^{\prime}} \underbrace{T^{\varrho^{\prime}} S_{B}^{\varrho, \sigma}\left(T^{\sigma^{\prime}}\right)^{\top}}_{=: S_{B}^{\rho_{B}^{\prime}, \sigma^{\prime}}}\left(V^{\sigma^{\prime}}\right)^{\top} .
$$

Complexity: Only $\mathcal{O}\left(C_{\text {sp }}^{2} n\right)$ such splitting operations are required $\Rightarrow$ Complexity $\mathcal{O}\left(C_{\mathrm{sp}}^{2} n k^{3}\right)$.
Result: The best approximation $C$ of $A B$ in a prescribed $\mathcal{H}^{2}$-matrix format can be computed in $\mathcal{O}\left(C_{\mathrm{sp}}^{2} n k^{3}\right)$ operations.

Example: Discretized double layer potential on the unit sphere $K$, approximate $K K$.

| $n$ | Build | $C_{\mathrm{sp}}$ | $\mathcal{H}^{2}$-MVM | $\mathcal{H}^{2}$-MMM | $\mathcal{H}$-MMM | Error |
| ---: | ---: | ---: | ---: | ---: | ---: | :--- |
| 512 | 1.7 | 16 | $<0.01$ | 0.3 | 0.7 | $1.1_{-2}$ |
| 2048 | 10.7 | 36 | 0.01 | 6.0 | 39.0 | $1.2_{-2}$ |
| 8192 | 49.4 | 42 | 0.10 | 35.6 | 315.9 | $4.5_{-3}$ |
| 32768 | 206.3 | 56 | 0.45 | 172.3 | 2085.2 | $4.9_{-3}$ |
| 131072 | 858.3 | 84 | 2.00 | 924.2 | 13213.8 | $8.6_{-5}$ |

Important: Here, the result matrix is projected to the same cluster bases and the same block structure as $K$.

Question: Can we reduce the error by using adapted cluster bases?

Modification: Choose cluster bases for result matrix adaptively to improve the precision.

| $n$ | $C_{\mathrm{sp}}$ | Old MMM | Error | New basis | New MMM | Error |
| ---: | ---: | ---: | :--- | ---: | ---: | ---: |
| 512 | 16 | 0.3 | $1.1_{-2}$ | 0.3 | 0.3 | $5.0_{-5}$ |
| 2048 | 36 | 6.0 | $1.2_{-2}$ | 9.5 | 6.8 | $6.6_{-5}$ |
| 8192 | 42 | 35.6 | $4.5_{-3}$ | 98.9 | 40.8 | $1.4_{-4}$ |
| 32768 | 56 | 172.3 | $4.9_{-3}$ | 878.9 | 194.7 | $4.3_{-6}$ |
| 131072 | 84 | 924.2 | $8.6_{-5}$ | 5030.2 | 966.9 | $2.99_{-6}$ |

Advantage: Precision significantly improved, computation time only slightly increased.

Current work: Improve the efficiency of the construction of adapted cluster bases.

Problem: For each admissble block $V^{\tau} S^{\tau, \sigma}\left(V^{\sigma}\right)^{\top}$ of the $\mathcal{H}^{2}$-matrix, the kernel function $\kappa$ has to be evaluated in $m^{2 d}$ points.
$\Rightarrow$ Very time-consuming if high precision is required.
Approach: Apply adaptive cross approximation to the coefficient matrices $S^{\tau, \sigma}$ to find rank $k$ representation $S^{\tau, \sigma} \approx X^{\tau, \sigma}\left(Y^{\tau, \sigma}\right)^{\top}$.

Result: Only $m^{d} k \ll m^{2 d}$ kernel evaluations required to approximate coefficient matrix $S^{\tau, \sigma}$.

Important: ACA is only applied to coefficient matrices.
$\Rightarrow$ Construction still linear in $n$, no problems with quadrature.

| $m$ | 5 | 6 | 7 | 8 |
| ---: | ---: | ---: | ---: | ---: |
| Std | $24.1 s$ | $64.7 s$ | $160.8 s$ | $399.2 s$ |
| HCA | $22.3 s$ | $49.2 s$ | $90.0 s$ | $222.7 s$ |

$\mathcal{H}^{2}$-matrices: Data-sparse approximation of certain dense matrices with linear complexity in the number of degrees of freedom.

Adaptive cluster bases: Find quasi-optimal expansion systems for arbitrary matrices. In the case of integral operators, point evaluations of the kernel function are sufficient to construct efficient and reliable $\mathcal{H}^{2}$-matrix approximations.
$\mathcal{H}^{2}$-matrix arithmetics: Approximate $A+\lambda B$ and $A B$ in $\mathcal{O}\left(n k^{2}\right)$ and $\mathcal{O}\left(n k^{3}\right)$ operations. Next goal: Approximate $A^{-1}$ and $L U$.

HCA: Reduce the number of kernel evaluations required to construct the $\mathcal{H}^{2}$-matrix.

Software: HLIB package for hierarchical matrices and $\mathcal{H}^{2}$-matrices, including applications to elliptic PDEs and integral equations. http://www.hlib.org


Construction: Pointwise evaluation and integrals of polynomials.

$$
T_{\nu^{\prime} \nu}^{t^{\prime}}=\mathcal{L}_{\nu}^{t}\left(x_{\nu^{\prime}}^{t^{\prime}}\right), \quad S_{\nu \mu}^{t, s}=\kappa\left(x_{\nu}^{t}, x_{\mu}^{s}\right), \quad V_{i \nu}^{t}=\int_{\Gamma} \varphi_{i}(x) \mathcal{L}_{\nu}^{t}(x) d x
$$

| Interpolation | Multipole |
| :--- | :--- |
| - Rank $k=p^{d}$ | + Rank $k=p^{d-1}$ |
| - Low frequencies only | + Low and high frequencies |
| + General kernels | - Special kernels only |
| + Anisotropic clusters | - Spherical clusters only |



Goal: Reduce rank for interpolation without sacrificing precision.

