

Hamiltonicity: sufficient conditions (contd.)

Theorem 4.

(Chvátal, Erdős 1972)

Consider a graph G with $n := |G| \geq 3$. The inequality $\kappa(G) \geq \alpha(G)$ implies the hamiltonicity of G .

($\alpha(G)$ is the stability number of G , i.e.

$\alpha(G) := \max\{|S| : S \subseteq V(G) \text{ and } G[S] \text{ has no edges}\}.$)

The result of Theorem 4 is best possible:

- (i) There exists a graph G with $\kappa(G) = \alpha(G) - 1$ which is not Hamiltonian (e.g. the Petersen graph or $K_{r,r+1}$, as discussed in the practical)
- (ii) The statement also does not hold if $\kappa(G)$ is replaced by $\lambda(G)$.

Corollary 5.

There exists an $O(mn)$ algorithm which either constructs a Hamiltonian cycle in G (with $n := |V(G)|$, $m := |E(G)|$) or finds a stable set S and a separating set T in G such that $|T| < |S|$ holds.

Hamiltonicity: further sufficient conditions

Theorem 6.

(Bondy, 1978)

Let G be a graph with $|G| \geq 3$ and $\deg(u) + \deg(v) \geq |G|$ for all $u, v \in V(G)$ with $\{u, v\} \notin E(G)$. Then $\kappa(G) \geq \alpha(G)$ holds.

Finding long cycles in graphs

The problem of determining the length of a longest cycle in a graph G is *NP*-hard.

We assume w.l.o.g. that G is 2-connected.

Theorem 7.

(Ore 1960, Bermond 1976, Linial 1976)

Let $d \in \mathbb{N}$ and let G be a 2-connected graph such that $\forall u, v \in V(G)$ with $\{u, v\} \notin E(G)$, $\deg(u) + \deg(v) \geq d$ holds. Then there exists a cycle of length at least $\min\{n, d\}$ in G , where $n := |G|$.

Theorem 8.

(Bauer, Broersma, Veldman and Rao 1989)

If G is a 2-connected graph with n vertices and connectivity $\kappa(G)$ such that $\deg(x) + \deg(y) + \deg(z) \geq n + \kappa(G)$ for any triple of independent vertices $x, y, z \in V(G)$, then G is Hamiltonian.

Proof in *Journal of Combinatorial Theory, Series B* **47(2)**, 1989, 237–243.

Hamiltonian cycles and degree sequences

Definition 1.

If G is a graph with $n := |G|$ and vertex degrees $d_1 \leq d_2 \leq \dots \leq d_n$, then the n -tuple (d_1, d_2, \dots, d_n) is called **the degree sequence** of G . An arbitrary sequence of natural numbers (a_1, a_2, \dots, a_n) is called a **Hamiltonian sequence** if every graph with n vertices and degree sequence (d_1, d_2, \dots, d_n) which is pointwise not smaller than (a_1, a_2, \dots, a_n) (i.e. $d_i \geq a_i, \forall i \in 1, n$), is Hamiltonian.

The degree sequence of a graph is unique while there may be different enumerations of the vertices all of them leading to the same degree sequence.

Theorem 9.

(Chvátal 1972)

A sequence (a_1, a_2, \dots, a_n) of natural numbers with

$0 < a_1 \leq a_2 \leq \dots \leq a_n < n$ and $n \geq 3$ is Hamiltonian iff $a_i \leq i$ implies $a_{n-i} \geq n - i$, for every $i < \frac{n}{2}$.