

# Claims involved in the proof of Edmonds-Gallai Decomposition Theorem

Consider the maximum cardinality matching produced by Edmonds' blossom algorithm applied to an undirected graph  $G = (V, E)$ . Let  $X$  be the set of vertices which are unmatched in  $M$ . Define the following three subsets of  $V$ :

$$\text{Even} := \{v \in V : \exists x \in X \text{ and an } M\text{-alternating path of even length from } x \text{ to } v\}$$

$$\text{Odd} := \{v \in V : \exists x \in X \text{ and an } M\text{-alternating path from } x \text{ to } v\} \setminus \text{Even}$$

$$\text{Free} := \{v \in V : \forall x \in X \text{ there exists no } M\text{-alternating path from } x \text{ to } v\}$$

Let *shrunk graph*  $G_0$  be the graph obtained in the final iteration of Edmonds' algorithm on  $G$ .

**Claim 1** If  $(u, v) \in E$  and  $u \in \text{Even}$ , then there is an  $M$ -alternating walk of odd length from  $X$  to  $v$  and there is an  $M$ -alternating path from  $X$  to  $v$ .

**Corollary** There is no edge between Even and Free in  $G$ .

**Claim 2** In  $G_0$  there is no edge between two even vertices.

**Claim 3**  $\text{Even} = D(G) := \{v \in V : \text{there exists a maximum cardinality matching in } G \text{ which does not match } v\}$

**Claim 4**  $\text{Odd} = A(G) := \{v \in V : v \text{ is neighbor of some } u \in D(G) \text{ but } v \notin D(G)\}$

**Claim 5**  $\text{Free} = C(G) := V \setminus (D(G) \cup A(G))$

**Claim 6**  $|M \cap C(G)| := \frac{|C(G)|}{2}$

**Claim 7** For every connected component  $H$  of  $(G \setminus A(G)) \cap D(G)$ :

- (a) either  $|X \cap H| = 1$  and  $M \cap \delta(H) = \emptyset$  or  $X \cap H = \emptyset$  and  $|M \cap \delta(H)| = 1$ , where  $\delta(H)$  is the set of edges with exactly one endvertex in  $H$ .
- (b)  $H$  is factor-critical.

**Claim 8**  $|M| = \frac{1}{2}(|V| + |A(G)| - o(G \setminus A(G)))$