Ex. 1

- Let G be a graph with n vertices $(n \ge 2)$.
 - (i) Show that at least two distinct vertices u, v of G have the same valency $\rho(v)$. (See Ex 1 Q3(ii)(e).)
 - (ii) Describe those graphs G for which every vertex v has valency $\rho(v) = 1$.
- (iii) Prove that if $\rho(v) = 2$ for all the vertices v of a connected graph G, then G is isomorphic to C_n .

Ex. 2

Let G be a graph with n vertices $(n \ge 2)$ and m edges.

- (i) Show that for a bipartite graph $G, m \leq \frac{n^2}{4}$. Give an example with equality when n is even and find a similar upper bound and example for when n is odd.
- (ii) Show that if G has an isthmus e, then it has at least two odd vertices. (Hint: Assume that G does not have any odd vertices and apply Corollary 1.3.8 to the components of $G^* = G^*(V, E \setminus \{e\})$.
- (iii) Show that if the graph G is disconnected, then its complement \overline{G} defined in Exercises 2 Question 3 is a connected graph. What is the diameter of \overline{G} in this case?

Ex. 3

- (i) Prove that if ρ ≥ 2 for all the vertices v of a graph G, then G contains a cycle. (Notes: This can be deduced from the Handshaking Lemma 1.3.7 and Theorem 1.6.2, as used in the proof of Theorem 1.6.5), but try to find a direct proof. Although they can be proved in a similar way, this is not quite the same theorem as in Question 1(iii), though that result can easily be deduced from this one.)
- (ii) Deduce from (i) that a tree with at least 2 vertices has at least one vertex v with valency $\rho(v) = 1$.
- (iii) Use (ii) to give a simpler proof of Theorem 1.6.2. using weak induction on n. Deduce from this and Theorem 1.3.7 that a tree with at least 2 vertices has at least two vertices v with valency $\rho(v) = 1$. Note: The first part of (iii) would be meaningless if (i) had been deduced from Theorems 1 and 5 as there would then be a circular argument i.e. $P \longrightarrow Q \longrightarrow R \longrightarrow P$. There are

and 5, as there would then be a circular argument, i.e. $P \Longrightarrow Q \Longrightarrow R \Longrightarrow P$. There are usually several ways of ordering the proofs of these theorems.

(iv) Prove that every tree with 2 vertices is a bipartite graph.

Ex. 4

- (i) Let A, B, C, D be four distinct vertices. List all the labelled trees T(V, E) that can be formed from $V = \{A, B, C, D\}$. For this example the edge AB is not equivalent to AC or AD etc., so that the Line Graph L_4 with vertices labelled ABCD in that order is not the same as that labelled CABD, though it is the same as the one with vertices DCBA. So there are many different labelled trees which are isomorphic graphs as defined in the lectures.
- (ii) How many labelled trees as defined in (i) are there for a set of five vertices? Hint: There are far too many of them to list so try to count the numbers in each isomorphism class. For example a lot of the labelled trees will be isomorphic to the unlabelled tree L_5 .