

On cluster primes

by

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1. Introduction. Blecksmith, Erdős and Selfridge [1] defined a prime $p > 2$ to be a *cluster prime* if every positive even integer $2r \leq p - 3$ can be written as a difference of two primes, $2r = q - q'$, where $q' \leq q \leq p$. It is an open question whether there exist infinitely many cluster primes. Guy ([4, Section C1]) attributes this question to Erdős. The attention of the general audience was drawn to this problem by Peterson's article [6] in Science News.

Blecksmith *et al.* [1] proved that the counting function $\pi_C(x)$ of cluster primes can be bounded from above: for all positive s ,

$$\pi_C(x) = O_s\left(\frac{x}{(\log x)^s}\right).$$

It is the purpose of this note to prove a better bound, i.e. that cluster primes are rare. This new bound was indeed conjectured by Blecksmith *et al.* [1].

THEOREM. *The number $\pi_C(x)$ of cluster primes below x is bounded by*

$$\pi_C(x) = O\left(\frac{x}{\exp\left(\frac{1}{60}(\log \log x)^2\right)}\right).$$

As Blecksmith, Erdős and Selfridge show, the problem is related to the prime k -tuple conjecture. It is proved that for a cluster prime p the interval $[p - t, p)$ must contain sufficiently many primes, which explains the name cluster prime. This allows us to apply an upper bound sieve. In Blecksmith *et al.* [1], Brun's version of the small sieve is used. The principal problem is that the authors arrive at a constant M whose dependence on the sieve dimension s is not at all clear. This prohibits taking an increasing s .

Filaseta [3] mentioned that an application of Hooley's almost pure sieve proves the result with $s = \varepsilon \log \log \log x$, thus obtaining an upper bound of

$$\pi_C(x) = O\left(\frac{x}{\exp(a \log \log x \log \log \log x)}\right) \quad \text{for some positive constant } a.$$

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In this note we apply the large sieve method, due to Montgomery [5]. In fact, we make use of the following lemma due to Vaughan [7], which is an elaborated version of the large sieve method, perfectly fitting to our application.

LEMMA 1 (Montgomery [5], Vaughan [7]). *Denote by \mathcal{P} the set of primes and let $\omega : \mathcal{P} \rightarrow \mathbb{N}$ with $0 \leq \omega(p) \leq p - 1$. Let $\mathcal{A} \subset [1, N]$ denote a set of integers which lies outside $\omega(p)$ residue classes modulo the prime p . Then the number $A(x)$ of elements $n \in \mathcal{A}$ with $n \leq x$ satisfies*

$$A(x) \leq \frac{2x}{L}, \quad \text{where } L = \sum_{q \leq x^{1/2}} \mu^2(q) \prod_{p|q} \frac{\omega(p)}{p - \omega(p)}.$$

Moreover,

$$L \geq \max_{m \in \mathbb{N}} \exp \left(m \log \left(\frac{1}{m} \sum_{p \leq x^{1/(2m)}} \frac{\omega(p)}{p} \right) \right).$$

2. Proof of the Theorem. If p is a cluster prime, then the even integers like $p - 9$ or $p - 15$ are the differences of two primes q, q' with $q, q' \leq p$. In particular, there must be a prime in the interval $[p - 6, p)$. More generally, an even integer $2r \in [p - t, p - 3]$ must be represented by a prime $q \in [p - t, p]$ and a prime $q' \in [1, t]$. By the prime number theorem the number of primes in $[1, t]$ is $(1 + o(1))t/\log t$. We see that for any $\varepsilon > 0$ there must be at least $s := (1/2 - \varepsilon) \log t$ primes in $[p - t, p)$.

Since the average gap between primes of size x is about $\log x$ we see that this is a useful criterion for $t = O((\log x)^\delta)$ (with $0 < \delta < 1$). On the contrary, for sufficiently large t one expects that an interval of length t has about $t/\log t$ primes so that this criterion becomes useless.

There are (trivially) at most $\binom{t}{s}$ possibilities to place s primes in an interval of length t . For any pattern of s primes in $[p - t, p)$ we will give an upper bound on the number of prime s -tuples below x . This bound will not depend on the particular pattern. So, multiplying this bound by the upper bound for the number of patterns, $\binom{t}{s}$, gives an upper bound on the number of $p \leq x$ such that the interval $[p - t, p)$ contains (at least) s primes.

We prove the following lemma:

LEMMA 2. *Let $\delta = 1/(2e)$ and $t = (\log x)^\delta$. Let ε be a sufficiently small positive constant. Let $A(x)$ denote the number of integers $n \leq x$ such that the interval $[n - t, n)$ contains at least $s = (1/2 - \varepsilon) \log t$ primes. Then*

$$A(x) = O \left(\frac{x}{\exp((1/(4e^2) - \varepsilon)(\log \log x)^2)} \right).$$

We fix a particular pattern $a_1 < \dots < a_s$. If all $n - a_i$ are prime simultaneously, then the integers n avoid the residue classes $a_i \pmod p$ for $p \leq n - t$. If $p > t$, then the number of forbidden classes is $\omega(p) = s$.

We choose $m = \lceil \delta^2(\log \log x)^2 \rceil$, where x is large. With $y = x^{1/(2m)}$ we have $\log \log y = \log \log x - 2 \log \log \log x + O(1)$. So we find that

$$\begin{aligned} \sum_{p \leq y} \frac{\omega(p)}{p} &\geq \sum_{t \leq p \leq y} \frac{\omega(p)}{p} \geq s(\log \log y - \log \log t + o(1)) \\ &\geq (1/2 - \varepsilon)\delta(\log \log x)(\log \log x - 3 \log \log \log x + O(1)) \\ &\geq (1/2 - 2\varepsilon)\delta(\log \log x)^2. \end{aligned}$$

This implies the estimate

$$\begin{aligned} L &\geq \exp \left(\lceil \delta^2(\log \log x)^2 \rceil \log \left(\frac{1}{\lceil \delta^2(\log \log x)^2 \rceil} \delta \left(\frac{1}{2} - 2\varepsilon \right) (\log \log x)^2 \right) \right) \\ &\geq \exp \left(\delta^2(\log \log x)^2 \log \left(\frac{1}{\delta} \left(\frac{1}{2} - 3\varepsilon \right) \right) \right) \\ &\geq \exp(\delta^2(\log \log x)^2 \log(e - 6e\varepsilon)) \geq \exp(\delta^2(1 - 7\varepsilon)(\log \log x)^2) \\ &\geq \exp \left(\left(\frac{1}{4e^2} - \varepsilon \right) (\log \log x)^2 \right). \end{aligned}$$

Therefore, for any fixed pattern $a_1 < \dots < a_s$ there are at most

$$\frac{2x}{\exp((1/(4e^2) - \varepsilon)(\log \log x)^2)}$$

values $n \leq x$ such that all $n - a_i$ are prime. Thus the lemma is proved.

To prove the theorem we only need to recall that

$$\pi_C(x) \leq \binom{t}{s} \frac{2x}{\exp(\delta^2(1 - 7\varepsilon)(\log \log x)^2)}.$$

Because of

$$\binom{t}{s} \leq t^s \leq \exp((1/2 - \varepsilon)\delta^2(\log \log x)^2)$$

we find that

$$\pi_C(x) = O \left(\frac{x}{\exp((1/(8e^2) - \varepsilon)(\log \log x)^2)} \right).$$

3. Further comments. No serious attempt has been made at optimizing the constant $1/60$ or $1/(8e^2) - \varepsilon$ that appears in the Theorem. Some improvement is possible. We only mention the following: Vaughan's argument in Lemma 1 can be refined to

$$L \geq \max_{m \in \mathbb{N}} \exp \left(m \log \left(\frac{e - \varepsilon_m}{m} \sum_{p \leq x^{1/(2m)}} \frac{\omega(p)}{p} \right) \right).$$

Here the ε_m are positive constants that tend to 0 as m goes to infinity. This allows using $c_1 \approx \delta/2$ and $\delta \approx 1/2$ and proves the Theorem with $1/8 - \varepsilon$ instead of $1/60$. For details see [2].

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